"Calibeating": Beating Forecasters at Their Own Game

Sergiu Hart

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Joint work with

Dean P. Foster

University of Pennsylvania & Amazon Research NY

Sergiu Hart "Calibration: The Minimax Proof", 1995 [2021]

www.ma.huji.ac.il/hart/publ.html#calib-minmax

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www.ma.huji.ac.il/hart/publ.html#calib-int

Dean P. Foster and Sergiu Hart "Forecast Hedging and Calibration" Journal of Political Economy 2021

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www.ma.huji.ac.il/hart/publ.html#calib-int
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Dean P. Foster and Sergiu Hart "'Calibeating': Beating Forecasters at Their Own Game" Theoretical Economics 2023

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www.ma.huji.ac.il/hart/publ.html#calib-beat
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Forecaster says: "The probability of rain tomorrow is p"

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- Forecaster is CALIBRATED if
 - for every forecast p: in the days when the forecast was p, the proportion of rainy days equals p(or: is close to p in the long run)

CALIBRATION can be guaranteed

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(no matter what the weather will be)

Foster and Vohra 1994 [publ 1998]

CALIBRATION can be guaranteed

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CALIBRATION can be guaranteed

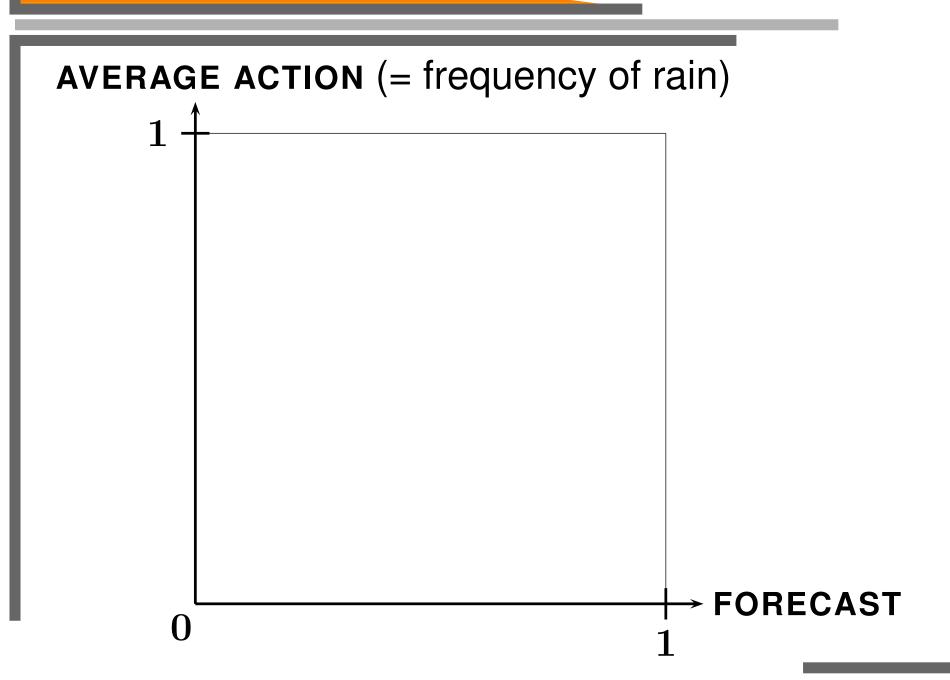
- Foster and Vohra 1994 [publ 1998]
- Hart 1995: proof by Minimax Theorem
- **.** . . .
- Hart and Mas-Colell 1996 [publ 2000]: procedure by Blackwell's Approachability

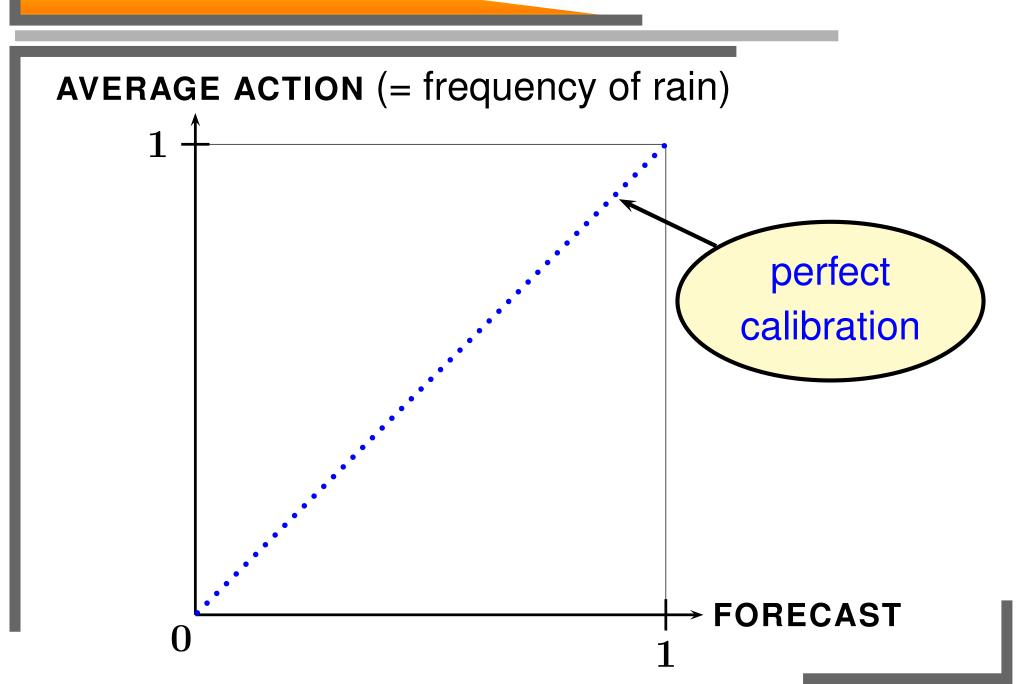
CALIBRATION can be guaranteed

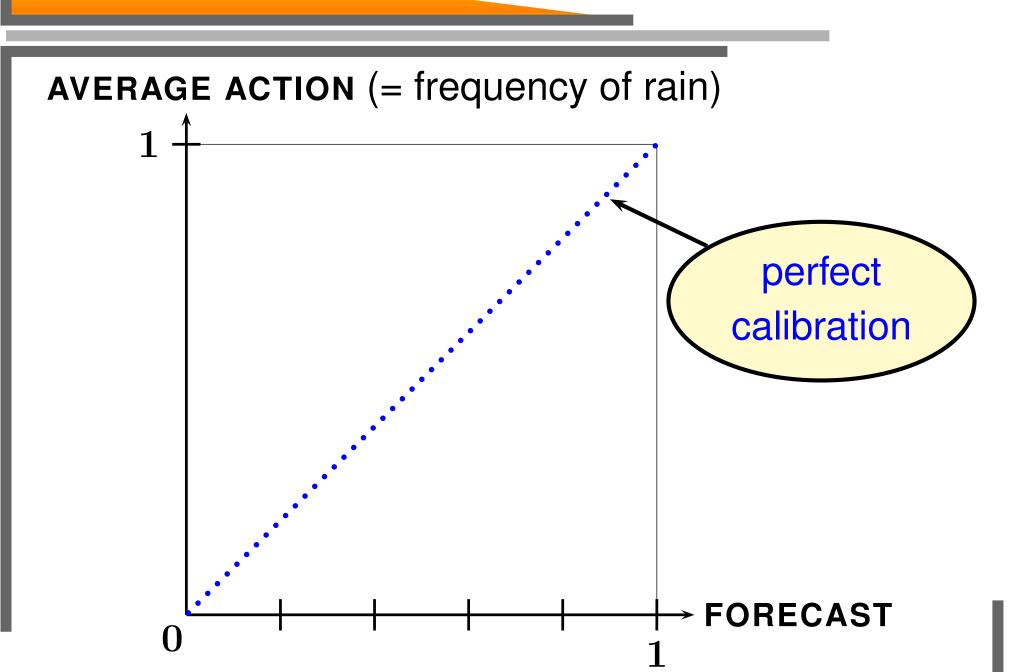
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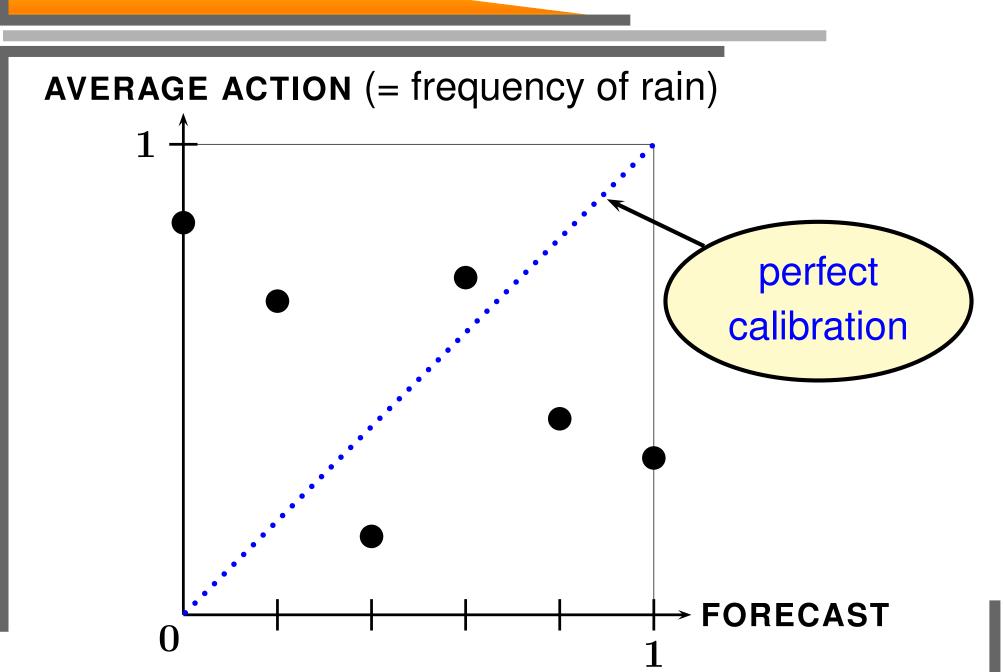
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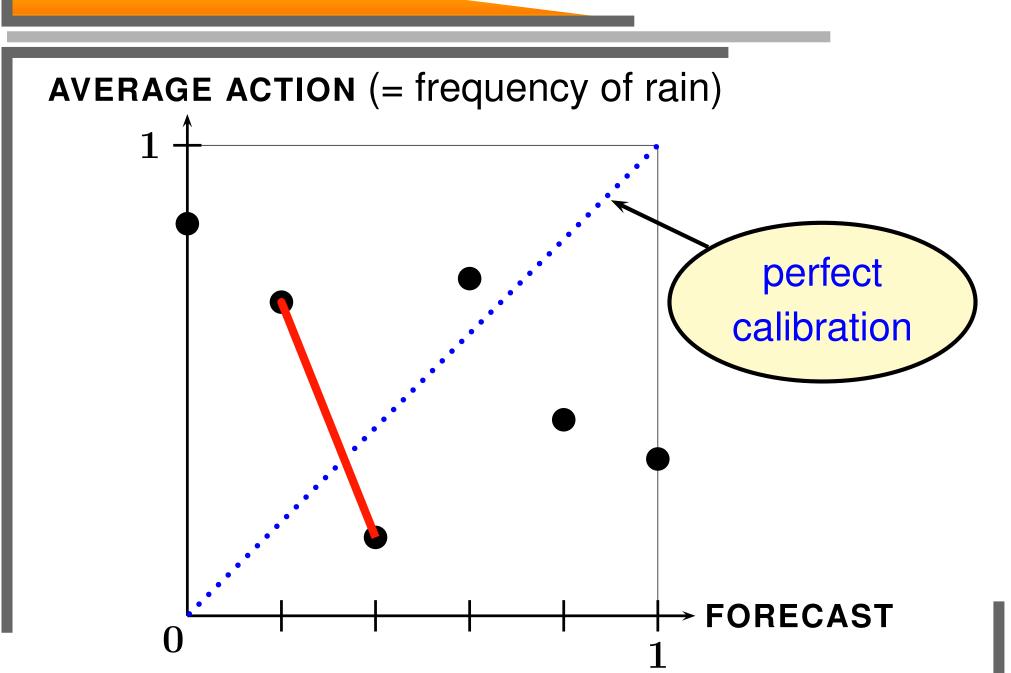
- Foster and Vohra 1994 [publ 1998]
- Hart 1995: proof by Minimax Theorem
- Hart and Mas-Colell 1996 [publ 2000]: procedure by Blackwell's Approachability
- Foster 1999: simple procedure
- Foster and Hart 2016 [publ 2021]: simplest procedure, by "Forecast Hedging"

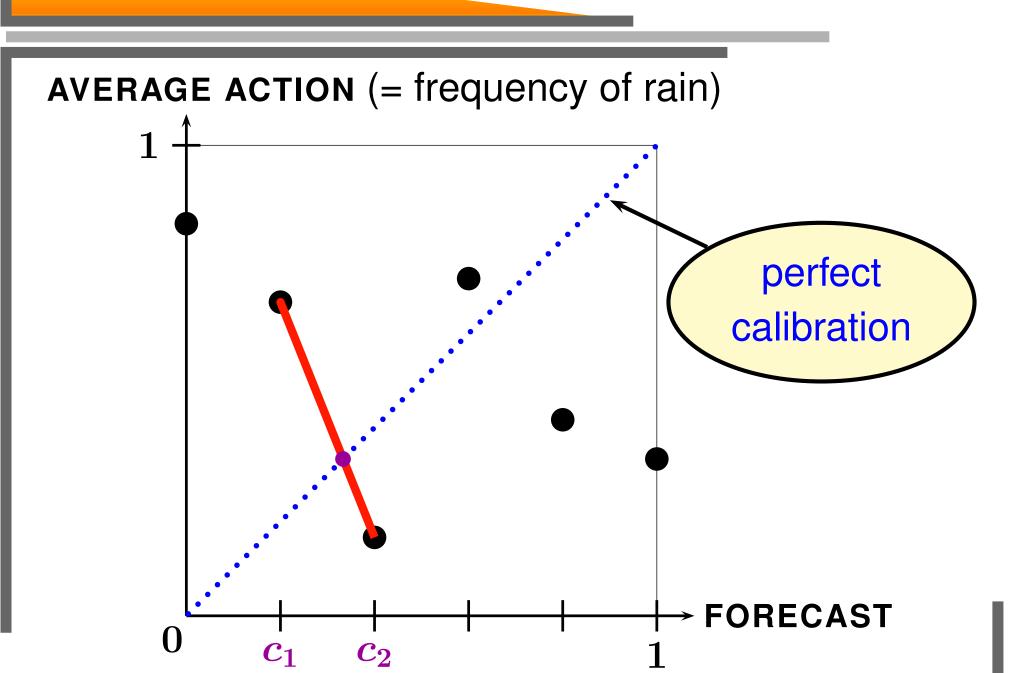






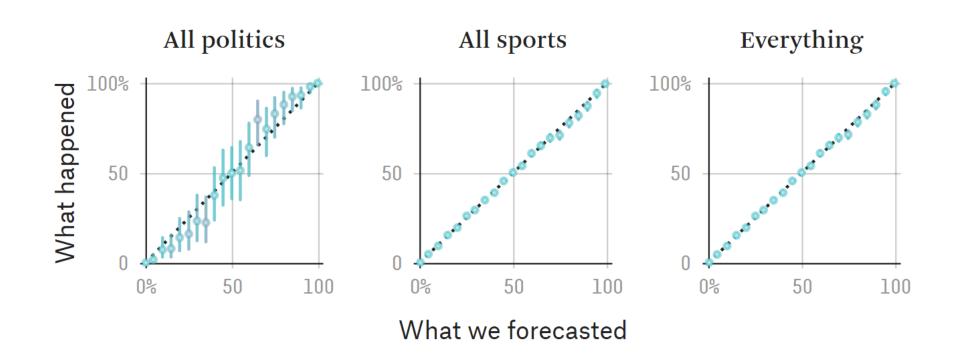






Calibration in Practice

Calibration in Practice



Calibration plots of FiveThirtyEight.com (as of June 2019)

Calibration in Practice



Prediction buckets

Calibration plot of ElectionBettingOdds.com (2016 – 2018)

time	1	2	3	$oxed{4}$	5	6	
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time	1	2	3	4	5	6	•••
rain	1	0	1	0	1	0	,

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F1: CALIBRATION = 0

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F1	100%	0%	100%	0%	100%	0%	
F2	50%	50%	50%	50%	50%	50%	

F1: CALIBRATION = 0

F2: CALIBRATION = 0

time	1	2	3	4	5	6	
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F1	100%	0%	100%	0%	100%	0%	
F2	50%	50%	50%	50%	50%	50%	

F1: CALIBRATION = 0 IN-BIN VARIANCE = 0

F2: CALIBRATION = 0

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F1: CALIBRATION = 0 IN-BIN VARIANCE = 0

F2: CALIBRATION = 0 IN-BIN VARIANCE = $\frac{1}{4}$

• $a_t = action at time t$

- $m{a}_t = ext{action at time } t$
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$$ar{a}(x) = rac{\sum_{t=1}^T \mathbf{1}_x(oldsymbol{c}_t)\,a_t}{\sum_{t=1}^T \mathbf{1}_x(oldsymbol{c}_t)}$$

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$$\mathcal{R} = rac{1}{T}\sum_{t=1}^T \|a_t - ar{a}(oldsymbol{c}_t)\|^2$$

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BRIER = REFINEMENT + CALIBRATION

Proof.

$$\mathbb{E}[(X-c)^2] = \mathbb{V}ar(X) + (ar{X}-c)^2$$

where c is a constant and $oldsymbol{X}$ is a random variable with $ar{oldsymbol{X}} = \mathbb{E}[oldsymbol{X}]$

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F1:
$$\mathcal{K} = 0$$
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F2:
$$K = 0$$
 $R = \frac{1}{4}$

time	1	2	3	4	5	6	
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F1:
$$\mathcal{K} = 0$$
 $\mathcal{R} = 0$ $\mathcal{B} = 0$

F2:
$$K = 0$$
 $R = \frac{1}{4}$ $B = \frac{1}{4}$



"Experts"

Testing experts:

"Experts"

Testing experts:

✓ BRIER score

"Experts"

Testing experts:

- **✓ Brier** score
- X CALIBRATION SCORE



"Expertise"

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LOW REFINEMENT SCORE

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- But: CALIBRATION procedures ignore whatever "EXPERTISE" one has

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Question:

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without LOSING "EXPERTISE"?

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ullet Can one get $\mathcal K$ to 0 without increasing $\mathcal R$?

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- But: CALIBRATION procedures ignore whatever "EXPERTISE" one has

Question:

Can one GAIN CALIBRATION without LOSING "EXPERTISE"?

- Can one get K to 0 without increasing R?
- Can one decrease $\mathcal{B} = \mathcal{R} + \mathcal{K}$ by \mathcal{K} ?

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 $\mathcal{K}^{'}=0$ $\mathcal{R}^{'}=\mathcal{R}$ $\mathcal{B}^{'}=\mathcal{B}-\mathcal{K}$

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• IN RETROSPECT / OFFLINE (when the $\bar{a}(c)$ are known)

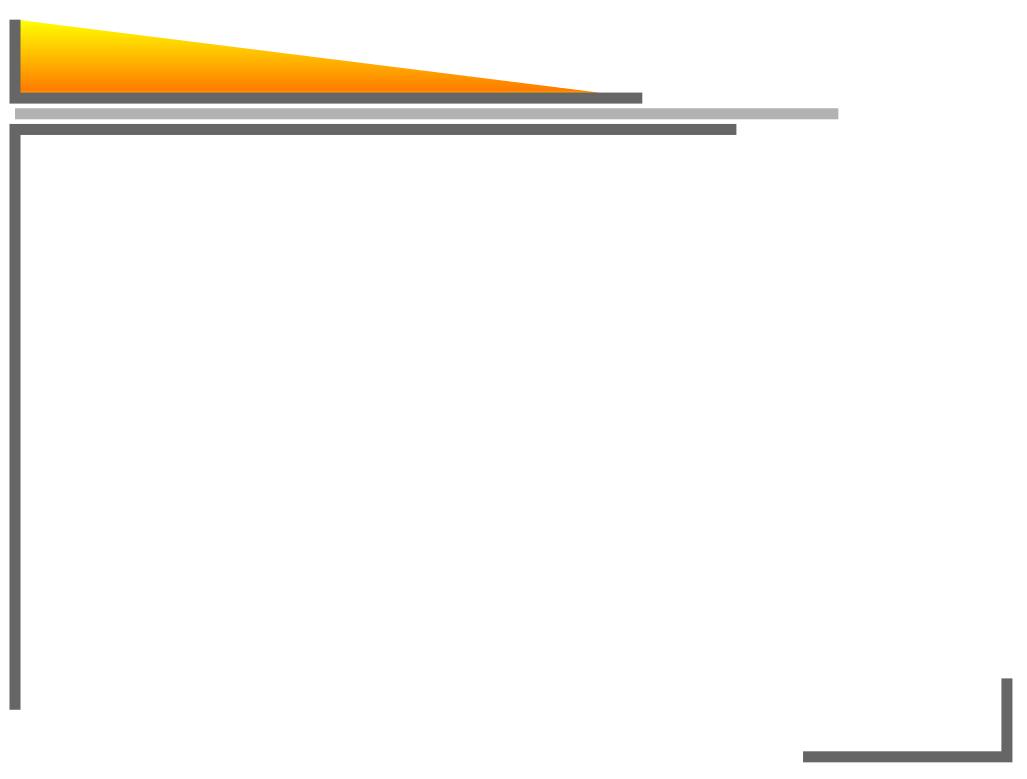
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Question:

Can one do this **ONLINE**?



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$$\mathcal{B}_T^{\mathrm{c}} \leq \mathcal{B}_T^{\mathrm{b}} - \mathcal{K}_T^{\mathrm{b}} + \mathrm{o}(1) \quad \mathrm{as} \ T \to \infty$$

for ALL sequences a_t and b_t (uniformly)

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 $oldsymbol{c}$ "BEATS" b by b 's CALIBRATION score

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c "BEATS" b by b's CALIBRATION score

GUARANTEED for ALL sequences of actions and forecasts

time	1	2	3	4	5	6	•••
rain	1	0	1	0	1	0	
b	80%	40%	80%	40%	80%	40%	

time	1	2	3	$oxed{4}$	5	6	•••
rain	1	0	1	0	1	0	
b	80%	40%	80%	40%	80%	40%	

b:
$$K^{\rm b} = 0.1$$
 $R^{\rm b} = 0$ $R^{\rm b} = 0.1$

$$\mathcal{R}^{\mathrm{b}} = 0$$

$${\cal B}^{
m b} = 0.1$$

time	1	2	3	4	5	6	
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b	80%	40%	80%	40%	80%	40%	
c	100%	0%	100%	0%	100%	0%	

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$$K^{b} = 0.1$$
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$$\mathcal{R}^{\mathrm{b}} = 0$$
 \mathcal{B}^{b}

$$c$$
: $\mathcal{K}^{c} = 0$

c:
$$\mathcal{K}^{c} = 0$$
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$$c$$
 calibeats b : $\mathcal{B}^{\mathrm{c}} \leq \mathcal{B}^{\mathrm{b}} - \mathcal{K}^{\mathrm{b}}$

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(that was easy ...)

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Can one CALIBEAT in general, non-stationary, situations?

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Weather is arbitrary and not stationary

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- Weather is arbitrary and not stationary
- Forecasts of b are arbitrary

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- **Binning of** b is not perfect ($\mathbb{R}^b > 0$)

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- Weather is arbitrary and not stationary
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- **Binning of** b is not perfect ($\mathbb{R}^b > 0$)
- Bin averages do not converge
- ONLINE
- GUARANTEED (even against adversary)

Theorem

There exists a **CALIBEATING** procedure

A Way to Calibeat

A Way to Calibeat

Theorem

The procedure

$$oldsymbol{c}_t = ar{a}_{t-1}^{ ext{b}}(b_t)$$

GUARANTEES b-CALIBEATING

A Simple Way to Calibeat

Theorem

The procedure

$$oldsymbol{c}_t = ar{a}_{t-1}^{ ext{b}}(b_t)$$

GUARANTEES b-CALIBEATING

Forecast the average action of the current *b*-forecast

$$\mathbb{V} ext{ar} \; = \; rac{1}{T} \sum_{t=1}^{T} \left\| x_t - ar{x}_T
ight\|^2$$

$$egin{array}{lll} \mathbb{V} \mathrm{ar} &=& rac{1}{T} \sum_{t=1}^{T} \left\| x_t - ar{x}_T
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(*)
$$o(1) = O\left(\frac{1}{T}\sum_{t=1}^{T} \frac{1}{t}\right) = O\left(\frac{\log T}{T}\right)$$

$$egin{array}{lll} \mathbb{V} \mathrm{ar} &=& rac{1}{T} \sum_{t=1}^{T} \|x_t - ar{x}_T\|^2 \ &=& rac{1}{T} \sum_{t=1}^{T} \left(1 - rac{1}{t}
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Proof: "Online Variance"

$$\mathbb{V}\text{ar} = \frac{1}{T} \sum_{t=1}^{T} \|x_t - \bar{x}_T\|^2 \\
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= \frac{1}{T} \sum_{t=1}^{T} \|x_t - \bar{x}_{t-1}\|^2 - o(1) \\
= \widetilde{\mathbb{V}\text{ar}} - o(1)$$

Proof: "Online Variance"

$$\mathbb{V}\mathrm{ar} = \mathbb{V}\mathrm{ar} - \mathrm{o}(1)$$

Proof: "Online Refinement"

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Theorem

$$\left[oldsymbol{c}_t = ar{a}_{t-1}^{\mathrm{b}}(b_t)
ight]$$

GUARANTEES b-CALIBEATING:

$$\mathcal{B}^{\mathrm{c}} \leq \mathcal{B}^{\mathrm{b}} - \mathcal{K}^{\mathrm{b}}$$

Self-Calibeating

Theorem

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Obtained by solving a Minimax problem (LP)

Outgoing Minimax (FH)

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- Obtained by solving a MINIMAX problem (LP)
- Moreover: solving a FIXED POINT problem yields a probability distribution P that is **ALMOST DETERMINISTIC**: its support is included in a ball of size δ

Theorem

There is a stochastic procedure that **GUARANTEES CALIBRATION**

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Proof. Self-calibeating + Stochastic Fixed Point

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Note. δ -CALIBRATION

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Proof. Calibeat the **joint** binning of b and c, by applying Stochastic Fixed Point

CONTINUOUS CALIBRATION: combine the days when the forecast was close to p

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Foster and Kakade (2004, 2006) Foster and Hart (2018, **2021**)

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Proof. Calibeat the **joint** binning of b and c, by a Fixed Point result (Brouwer)

Multi-Calibeating

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Theorem

There is a *deterministic* procedure that **GUARANTEES**

simultaneous CALIBEATING of several forecasters

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There is a **stochastic** procedure that **GUARANTEES**

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Proof. Calibeat the joint binning

In all the results above:

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	CALIBRATION	
Obtained by	Minimax	
Procedure	stochastic	

... and Continuous Calibration

In all the results above:

	CALIBRATION	CONTINUOUS
Obtained by	Minimax	Fixed Point
Procedure	stochastic	deterministic

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Proof. Label each bin b with $\phi(b) = \bar{a}_T(b)$

• c CALIBEATS b:

$$\mathcal{B}_T(\mathbf{c}) \leq \mathcal{R}_T(\mathbf{b}) + o(1)$$

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CALIBEATING is **stronger**!

TAKING PRIDE IN OUR RECORD

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"We have correctly forecasted 8 of the last 5 recessions"



TAKING PRIDE IN OUR RECORD

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